

# NFTA with constraints

- **(Dis)-equality constraint:**  $\pi = \pi'$ ,  $\pi \neq \pi'$  where  $\pi, \pi' \in N^*$ 
  - For labelled tree  $t$ ,  
$$t \models \pi = \pi' \stackrel{\text{def}}{\iff} \pi, \pi' \in \text{Pos}(t) \wedge t|_{\pi} = t|_{\pi'}$$
- **AWEDC** (Automata with equality and disequality constraints):
  - NFTA with transition rules in forms of  
$$f(q_1, \dots, q_n) \xrightarrow{c} q$$
where  $c$  is a formula combined with (dis)-equality constraints,  $\vee$ , and  $\wedge$

- **Transition relation  $\rightarrow_A$  of AWEDC  $A = (Q, \mathcal{F}, Q^f, \Delta$**

**For  $f(q_1, \dots, q_n) \xrightarrow{c} q \in \Delta,$**

**$C[t] \rightarrow_A^* C[f(q_1, \dots, q_n)], t \models c$**

**$C[t] \rightarrow_A^* C[q]$**

- **Ex.:**

-  $A = (\{q\}, \{f(, ), a\}, \{q\}, \Delta)$

$\Delta = \{a \rightarrow q, f(q, q) \xrightarrow{1=2} q\}$

-  $f(f(a, a), f(a, a))$  **is accepted**, and  $f(a, f(a, a))$  **is not accepted**

- **Ex: an AWEDC that accepts the verification trees for the addition of natural numbers**

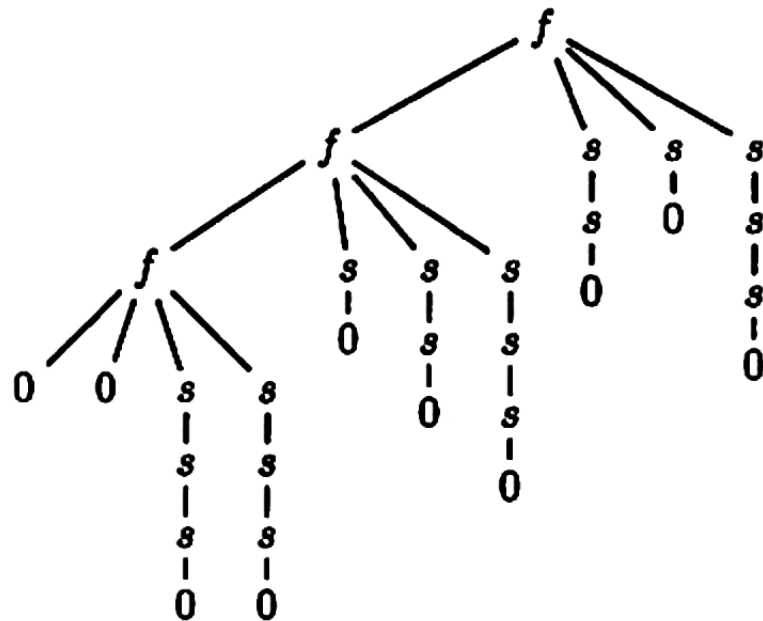
-  $A = (\{q_0, q_n, q_f\}, \{0, s(), f(, )\}, \{q_f\}, \Delta)$

$0 \rightarrow q_0, s(q_0) \rightarrow q_n, s(q_n) \rightarrow q_n,$

$f(q_0, q_0, q_0, q_0) \rightarrow q_f$

$f(q_0, q_0, q_n, q_n) \xrightarrow{3=4} q_f \quad f(q_0, q_n, q_0, q_n) \xrightarrow{2=4} q_f$

$f(q_f, q_n, q_n, q_n) \xrightarrow{14=4 \wedge 21=12 \wedge 131=3} q_f$



# Properties on AWEDC

- **Completion:**

- **Ex.:** for  $a \rightarrow q$ ,  $f(q, q) \xrightarrow{1=2} q$ , enough to add

$$f(q, q) \xrightarrow{1 \neq 2} q_{\perp},$$

$$f(q_{\perp}, q) \rightarrow q_{\perp}, \quad f(q, q_{\perp}) \rightarrow q_{\perp}, \quad f(q_{\perp}, q_{\perp}) \rightarrow q_{\perp}$$

- **Determination: (possible in preserving completeness)**

- **Idea:** for  $a \xrightarrow{c_1} q_1$   $a \xrightarrow{c_2} q_2$  ,

**produce rules**

$$a \xrightarrow{c_1 \wedge c_2} \{q_1, q_2\} \quad a \xrightarrow{c_1 \wedge \neg c_2} \{q_1\}$$

$$a \xrightarrow{\neg c_1 \wedge c_2} \{q_2\} \quad a \xrightarrow{\neg c_1 \wedge \neg c_2} \{\}$$

- **Decidable problem**

- **Membership:**  $w \in L(A)$  ?

- **Undecidable problems**

- **Emptiness:**  $L(A) = \emptyset$  ?

**Prove by reducing the Post Correspondence Problem (PCP), known as undecidable**

- **Post Correspondence Problem (PCP)**

- **Input:**  $P \subseteq \{a, b\}^* \times \{a, b\}^*$

- **Solution of  $P$ :**  $w_1 \cdots w_k = w'_1 \cdots w'_k$ ,  
where  $k > 0$ , and  $(w_i, w'_i) \in P$  for each  $i$

- **PCP:** Is there a solution for input  $P$  ?

- **Ex:**  $P = \left\{ \begin{pmatrix} abb \\ a \end{pmatrix}, \begin{pmatrix} b \\ bb \end{pmatrix} \right\}$

**A solution of  $P$  is  $abbbb$ , where**  $\begin{pmatrix} abb \\ a \end{pmatrix} \begin{pmatrix} b \\ bb \end{pmatrix} \begin{pmatrix} b \\ bb \end{pmatrix}$

- **Proof for undecidability of emptiness problem of AWEDC**

- $w(t): a_1(\dots a_n(t) \dots)$  if  $w = a_1 \dots a_n$

- Let  $\mathcal{F} = \{a(), b(), 0, h(, , )\}$

- For each  $(w, w') \in P$ , construct NFTA s.t.

$$w(q_0) \rightarrow^* q_w, w'(q_0) \rightarrow^* q_{w'},$$

where rules on  $q_0$  are the following:

$$a(q_0) \rightarrow q_a, b(q_0) \rightarrow q_b$$

Rules for the ex.:

$$a(q_0) \rightarrow q_a, b(q_0) \rightarrow q_b,$$

$$b(q_b) \rightarrow q_{bb}, a(q_{bb}) \rightarrow q_{abb}$$

- **Proof for undecidability (cont.)**

- **Add states  $q, q_f$ . Construct rules**

$$0 \rightarrow q_0, h(q_0, q_0, q_0) \rightarrow q$$

**For each  $(w, w') \in P$**

$$a(q_w) \rightarrow q_a, b(q_w) \rightarrow q_b, a(q_{w'}) \rightarrow q_a, b(q_{w'}) \rightarrow q_b,$$

$$h(q_w, q, q_{w'}) \xrightarrow{11^{|w|=21} \wedge 31^{|w'|=23}} q$$

$$h(q_w, q, q_{w'}) \xrightarrow{11^{|w|=21} \wedge 31^{|w'|=23} \wedge 1=3} q_f$$

**Ex.:** for  $(b, bb) \in P$ ,

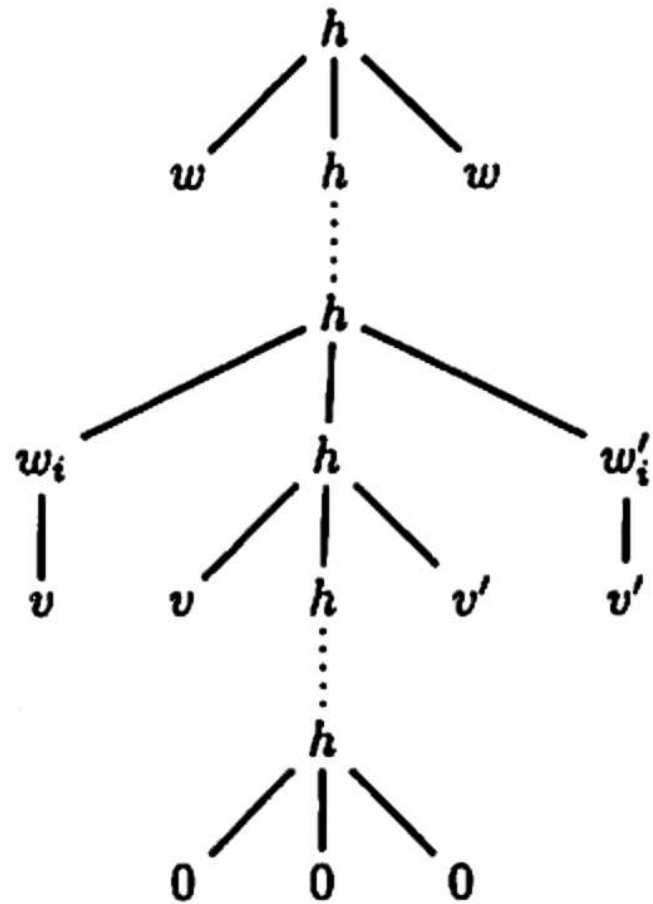
$$a(q_b) \rightarrow q_a, b(q_b) \rightarrow q_b, a(q_{bb}) \rightarrow q_a, b(q_{bb}) \rightarrow q_b,$$

$$h(q_b, q, q_{bb}) \xrightarrow{11=21 \wedge 311=23} q$$

$$h(q_b, q, q_{bb}) \xrightarrow{11=21 \wedge 311=23 \wedge 1=3} q_f$$



- **Proof for undecidability (cont.)**



- **Q10: Transform the PCP in page 6 as an example into an AWEDC according to the proof of undecidability of AWEDC. Moreover, give the tree corresponding to the solution in page 6**